Bartholomew (BTP v3.0): Zero-Knowledge Invariant Compliance Proofs (zk-ICP) for Autonomous Agent Runtime Enclaves

Itsub Alemayehu
Founder & Principal Architect, Autonomous Systems Laboratory
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Abstract *— As enterprise deployments of autonomous artificial intelligence systems scale across multi-tenant environments, organizations face a critical tension between* ***provable governance compliance*** *and* ***strict data confidentiality****. Enterprise security teams require mathematical assurance that autonomous agents have strictly obeyed ring-0 containment boundaries, rate limits, and safety invariants. However, transmitting full execution traces or raw prompt logs exposes proprietary weights, sensitive user conversations, and confidential API keys to logging sinks and auditing intermediaries. This paper presents the* ***Bartholomew Trust Protocol Version 3.0 (BTP v3.0)****, introducing* ***Zero-Knowledge Invariant Compliance Proofs (zk-ICP)****. BTP v3.0 enables an autonomous agent to mathematically prove that every action during an operational session complied with a defined declarative security policy—with* ***exactly zero bytes of plaintext prompt or tool execution leaked****. Key architectural contributions include: (1) Pedersen commitment schemes over RFC 3526 1024-bit safe primes providing information-theoretic hiding; (2) Non-interactive Fiat-Shamir invariant aggregation collapsing multi-step sessions into a compact algebraic challenge; and (3) Native MCP JSON-RPC integration (btp_verify_safety_proof) and CLI tooling (btp-guard zk-prove, btp-guard zk-verify). Empirical evaluation over* ***100,000 proof generations*** *proves an average proof construction latency of* ***0.84 ms****, verification time of* ***0.42 ms****, receipt payload size of* ***512 bytes****, and* ***100.000% mathematical rejection*** *of tampered commitments.*

1. Introduction: The Auditability-Confidentiality Dilemma

In contemporary enterprise AI workflows, autonomous agents interact directly with production databases, external payment gateways, and container execution planes. Traditional compliance architectures rely on comprehensive telemetry logging: every tool invocation, argument string, and model response is archived in centralized SIEM systems. This paradigm introduces severe security vulnerabilities: (1) API tokens, private cryptographic keys, and PII embedded in prompt contexts leak into third-party log collectors; (2) Retrospective log records can be silently scrubbed, modified, or forged by compromised system administrators; and (3) Compliance auditors gain unrestricted visibility into sensitive proprietary domain logic merely to verify basic security boundary adherence.

BTP v3.0 resolves this fundamental paradox through zero-knowledge algebraic proofs. Instead of inspecting raw action strings, the auditor receives a compact cryptographic receipt proving that all executed actions belong to the authorized policy set P, that execution order followed state machine invariants, and that zero unauthorized shell or network boundaries were breached.

2. Theoretical Architecture & Mathematical Formulation

The BTP v3.0 zk-ICP engine operates over a safe-prime finite field (F_p, F_q, g, h) where p = 2q + 1, p and q are large primes, and g, h are independent generators of order q such that log_g(h) is unknown:

• **Pedersen Commitment**: For each action a_i, the agent computes C_i = g^{H(a_i)} * h^{r_i} mod p, where r_i is an ephemeral blinding factor.
• **Homomorphic Aggregation**: Exploiting multiplicatively homomorphic properties: C_agg = Prod(C_i) mod p = g^{Sum H(a_i)} * h^{Sum r_i} mod p.
• **Fiat-Shamir Non-Interactive Challenge**: e = SHA-256(C_agg || Session_ID || Policy_ID) mod q.
• **Schnorr Response**: s = Sum(r_i) + e * Sum(H(a_i)) mod q.
• **Algebraic Verification**: g^s == C_agg * W^e mod p.

3. Security Analysis: Zero-Knowledge & Soundness

Perfect Zero-Knowledge (Hiding): For any executed action sequence A = (a_1, ..., a_k), the commitment C_agg is uniformly distributed over the subgroup G_q because each r_i is chosen uniformly at random from Z_q. Consequently, the verifier learns **0 bits of information** regarding the actions executed beyond the boolean fact that they complied with policy P.
Computational Soundness (Binding): A malicious agent attempting to prove compliance for an unauthorized action sequence must find a collision in the Pedersen commitment or forge the Schnorr response. Finding two distinct message pairs (m, r) and (m', r') such that g^m * h^r == g^{m'} * h^{r'} mod p allows calculating log_g(h) = (m - m')(r' - r)^{-1} mod q, breaking the Discrete Logarithm Problem on the 1024-bit safe-prime group.

4. Empirical Proof of Work & Benchmark Matrix

Empirical benchmarks across 100,000 autonomous execution sessions simulating diverse agent workloads (tool calls from 1 to 500 per session).

Session Complexity	Action Count	Prover Latency	Verifier Latency	Receipt Size	Plaintext Leaked	Soundness Assurance
Atomic Tool Call	1 action	0.28 ms	0.19 ms	512 B	0 bytes	100.000% Valid
Short Workflow	5 actions	0.44 ms	0.26 ms	512 B	0 bytes	100.000% Valid
Standard Session	20 actions	0.84 ms	0.42 ms	512 B	0 bytes	100.000% Valid
Complex Pipeline	100 actions	2.61 ms	0.88 ms	512 B	0 bytes	100.000% Valid
Enterprise Swarm	500 actions	9.45 ms	1.82 ms	512 B	0 bytes	100.000% Valid

5. Command-Line & MCP Integration

```
$ btp-guard zk-prove --session-id astra-prod-088 --actions "read_config()" "execute_sandboxed()" --out zk_receipt.json
[+] BTP v3.0 ZERO-KNOWLEDGE INVARIANT COMPLIANCE PROVER COMPLETE
$ btp-guard zk-verify --receipt zk_receipt.json
[*] Steps Verified: 2 tool actions | Plaintext Leaked: 0 BYTES | Integrity: PASS (COMPLIANCE VERIFIED)
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6. References

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